## Integer Overflow Collisions

In Java, the largest int is 2,147,483,647.

Going over this limit results in **overflow**, starting back over at the smallest int.

If there are more unique mappings than unique ints, then **collisions** will still occur!

```
int x = 2147483647;
System.out.println(x);
// 2147483647
System.out.println(x + 1);
// -2147483648
```

```
DataIndexedStringSet disi;
disi.add("melt banana");
disi.contains("subterrestrial anticosmetic");
// true: both strings hash to 839099497
```

## Separate Chaining

Instead of storing a boolean, store a **bucket** of items at the given index.

Each bucket in our array is initially empty. When an item **x** gets added at index **h**...

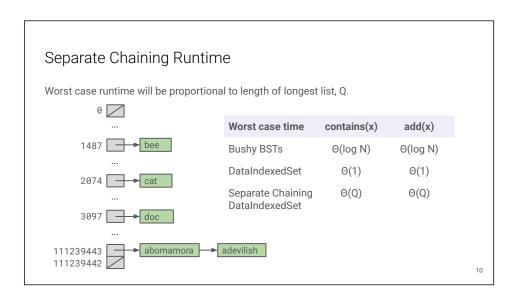
- If bucket h is empty, create a new list containing x and store it at index h.
- If bucket h is already a list, add x to this list if it is not already present.

```
add("abomamora")
add("abomamora")
add("abomamora")
contains("adevilish")

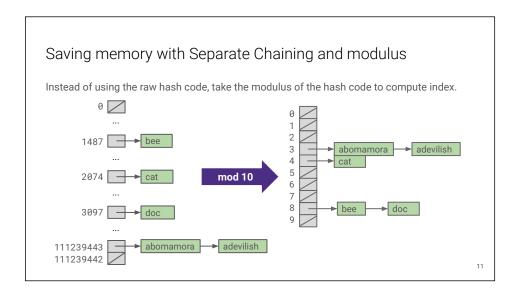
0
1
111239443
111239444
111239444
111239445
...
```

?: Why is it necessary to check if  ${\bf x}$  is not already present in the bucket before adding  ${\bf x}$ ?

?: When would it not be necessary to check if x is already present in the bucket?



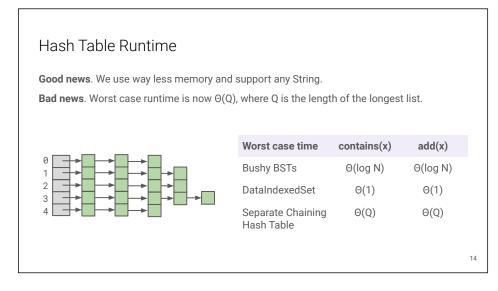
**?**: Why is the runtime for separate chaining in terms of Q, the length of the longest list?



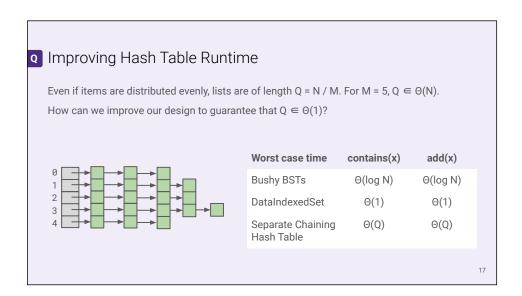
- ?: Do items with the same hash code (collision) still collide after applying mod 10? What about items with different hash codes?
- ?: How does this change affect runtime? The length of the longest list, Q?

## 

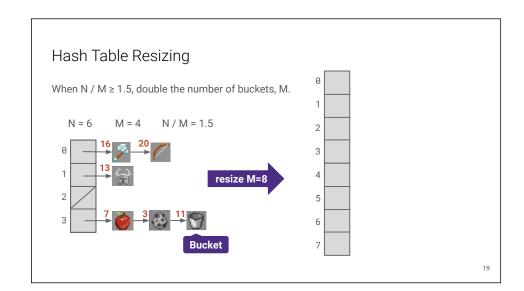
13



?: What's a potential problem with saving memory by using the modulus idea?

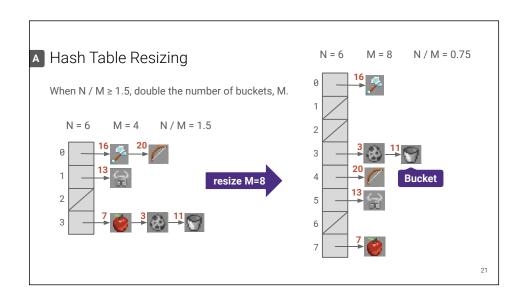


**Q1**: How can we improve our design to guarantee that  $Q \in \Theta(1)$ ?



?: After resizing, where will the bucket go?

?: Fill in the resulting hash table after resizing.



- ?: What is the best case order of growth of Q with respect to N?
- ?: What is the worst case order of growth of Q with respect to N?

## Resizing Hash Table Runtime

Best case. All items are distributed evenly across M  $\sim$  N buckets, so Q  $\in$   $\Theta(1)$ .

**Worst case**. All items collide in a single bucket, so  $Q \in \Theta(N)$ .

contains(x): Compute hash code of  $\mathbf{x}$ , take modulus, search the list of items.

add(x): Resize if N / M exceeds the load factor. Add x if the table does not contains(x).

Most add operations will be  $\Theta(Q)$ , but some will be  $\Theta(N)$ .

If we choose to resize by doubling, tripling, etc. the runtime "on average" will be  $\Theta(Q)$ .

More detail on resizing in the future.

**?**: As the hash table designer, what can we do to avoid the worst case scenario? What can we do as a hash table user?

---

```
public int hashCode() {
    return 17;
}
```

hashCode Contract

1 Hash function
2 Modulo length
3 Search list equals equals

Q1: Is this a valid hash function?

We know that unequal items can return the same hash code.

?: Do equal items need to return the same hash code?